Green and Sustainable Cloud of Things: Enabling Collaborative Edge Computing

Zhaolong Ning, Xiangjie Kong, Feng Xia, Weigang Hou, and Xiaojie Wang

ABSTRACT

The proliferation of IoTs beside the emergence of various cloud services push the horizon of edge computing. By offering cloud capabilities at the network edge closer to mobile devices, edge computing is a promising paradigm to resolve several vital challenges in IoTs, such as bandwidth saturation, energy constraints, low latency transmission, and data security and privacy. To provide a comprehensive understanding of edge computing supported by the integration of IoTs and cloud computing, that is, CoTs, this article first discusses some distinct research directions in CoTs with respect to edge computing. Given the significance of energy efficiency and sustainability of edge deployment in CoTs, we put forward a green and sustainable virtual network embedding framework for cooperative edge computing in wireless-optical broadband access networks. Specifically, we leverage a reliability function to confirm the number of backup edge devices, and embed virtual networks onto the suitable edge devices in CoTs. Finally, several research challenges and open issues are discussed.

INTRODUCTION

How to make the development of smart cities sustained is challenging due to the huge evolved investments. Owing to the superiorities of high computing capacity, competitive system scalability and low service cost, cloud computing has become a high-efficiency alternative in various fields of smart cities, such as social networks, commerce, transportation, education and health care services [1]. However, centralized cloud frameworks commonly face several challenges, such as network overload, end-to-end latency, single point of failure and high power consumption.

According to a report released by Cisco, mobile data traffic will increase by 53 percent annually from now until 2020 [2]. The number of network users provided by cloud services will be approximately 3.6 billion in 2018 (http://www.statista.com). With the advance triumphantly of Internet of Things (IoTs), the post-cloud era, that is, Cloud of Things (CoTs) is emerging. Thanks to the ever-increasing abilities of sensing, computing and communication of smart devices, nearly 50 billion things are likely to be connected through the Internet, where 45 percent of the generated data by IoTs will be handled by edge devices in 2019 [3]. However, it is overwhelmingly challenging for smart cities to manage or program such a huge system with massive mobile devices due to the high heterogeneity of IoT-based applications [4].

Currently, Wireless-Optical Broadband Access Networks (WOBAN) have become a key component in smart cities. As shown in Fig. 1, Wireless Mesh Networks (WMNs) and IoTs can be the front-end of WOBANs. The generated data are forwarded by routers via wireless channels and received by the neighbor optical network unit (ONU). Then, an optical line terminal (OLT) can receive the sensed data through a Passive Optical Network (PON). According to the obtained information flows, a decision can be made in the central office connecting to the OLT for industrial or commercial circumstances. With an ever-growing volume of data generated by smart devices and the saturating PON bandwidths, the central office becomes incapable of sustaining the continuous information flows. For the sake of alleviating the resource overprovisioning of a PON, deployment of edge devices (EDs) is advocated at both ONUs and wireless routers for collaborative edge computing, so that EDs can switch roles from data consumer to both data producer and data consumer.

Network virtualization technologies, such as Network Function Virtualization (NFV) and Software Defined Networking (SDN), are helpful to achieve highly-efficient edge deployment in CoTs by virtualizing device functions on top of standard hardware [5]. For example, a user’s computational capability can be abstracted to the Virtual Network Embedding (VNE) on the WMNs or IoTs infrastructure in WOBANs. Because of the inter-channel interference caused by EDs, VNE-based collaborative edge computing generally cannot guarantee network survivability. However, network survivability is important for sustainable edge computing in CoTs.

Although rapidly growing interest has been drawn to edge deployment, the related research of sustainable edge computing in CoTs is still in its infancy. To our best knowledge, our presented Green and Sustainable VNE (GSVNE) framework is the first effort to study the available number of backup EDs, and embed virtual networks (VNs) onto the suitable EDs for collaborative edge computing in CoTs. The main contributions of this article can be summarized as:

- We overview the current research studies of cooperative edge computing in CoTs, and classify the related studies into four categories.

The authors discuss some distinct research directions in CoTs with respect to the edge computing. Given the significance of energy efficiency and sustainability of edge deployment in CoTs, they put forward a green and sustainable virtual network embedding framework for cooperative edge computing in wireless-optical broadband access networks.

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The integration of fiber and wireless access networks in CoTs, and the determination of the number of EDs according to network reliability.

We select the qualified network by our cooperative scheme can both serve as many virtual networks as possible and provide a satisfied sharing degree of network backup resources.

The rest of this article is organized as follows. Some promising research directions for edge computing in CoTs are introduced in the next section. After that, the GSVNE framework is put forward. We then discuss research challenges and emerging solutions for collaborative edge computing in CoTs before concluding this article.

RESEARCH DIRECTIONS FOR EDGE COMPUTING IN COTS

Edge computing includes various technologies, such as cloudlet, local fog computing, sensing networks, cloud-centric IoTs and device-to-device (D2D) networks. In this section, we discuss the following four promising research directions for edge computing in CoTs.

EDGE COMPUTING OVER WOBANs

The integration of fiber and wireless access networks can offer the advantages of reliability in optical fiber networks and ubiquity in wireless networks. WOBANs can offer broadband services for both fixed subscribers and mobile clients. Therefore, it has become a powerful framework to provide widespread connectivity for services and applications in smart cities. Future smart cities not only require compatible radio interfaces, but also need the integration of heterogeneous networks to manage rapid data growth. Therefore, distributed edge computing over WOBANs is advocated. Furthermore, edge computing is also leveraged to satisfy the tactile-level latency requirements in future networks. This is because the total latency for mobile cloud computing is generally between 30 ms and 100 ms, which is intolerable for latency-critical applications, such as autonomous driving and online gaming [6].

Due to the heterogeneity and dense deployment of devices, edge caching is beneficial for boosting spectral efficiency and reducing energy consumption in CoTs. Content placement and delivery are two core issues in edge caching. The former considers the location, size, and downloaded content of network caches, while the latter focuses on delivering content to the requested terminals.

CACHE-ENABLED EDGE COMPUTING

Due to the heterogeneity and dense deployment of devices, edge caching is beneficial for boosting spectral efficiency and reducing energy consumption in CoTs. Content placement and delivery are two core issues in edge caching. The former considers the location, size, and downloaded content of network caches, while the latter focuses on delivering content to the requested terminals.

Service caching for resource allocation and data caching are two important challenges in this context. Although contents can be cached at WMNs or IoTs in WOBANs, the backhaul links are still constrained. The traffic in wired Internet, backhaul and mobile core networks can be reduced by deploying caches in base stations, relays and user devices [8]. Some trade-offs for edge caching deserve to be investigated, for example, the trade-offs between backhaul capacity and cache size, between throughput and outage in cache-enabled device-to-device (D2D) networks, and between energy consumption and offloading ratio.

COLLABORATIVE EDGE COMPUTING

Due to a local bottleneck, it is impossible for a single ED to cope with all computing tasks. Therefore, collaborative edge computing based on ED cooperation in the PON backhaul is promising to handle mission-critical applications. Also, collaborative edge computing is compatible with...
Since computation merely happens in the data facility of participants, data integrity and privacy can be well guaranteed. Some other promising topics include: role definition, resource mapping, computation offloading, data storage, and cooperation stimulation between cloud and edge entities.

<table>
<thead>
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<th>Definition</th>
<th>Description</th>
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<tr>
<td>$n$</td>
<td>Total number of WMNs and IoTs</td>
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<tr>
<td>$N$</td>
<td>Total number of VNs</td>
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<td>$N_1$</td>
<td>The number of embedded VNs processed by front-end WMNs and IoTs</td>
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<td>The number of VNs processed by the PON backhaul</td>
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<td>$\Gamma_m$</td>
<td>Working (backup) graph of the qualified WMN and IoT</td>
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<td>$\eta_i$</td>
<td>Binary variable to identify whether the $i$th VN is embedded onto $\Gamma_m$</td>
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Table 1. Main notations.

Virtualized Edge Computing

With the objective of enabling collaborative edge computing in CoTs for smart cities, flexible management of computing resources needs to be coordinated. Network virtualization is advocated, which is economical to customize computing tasks for mobile device users. The current network infrastructure largely benefits from the integration of NFV, SDN and edge computing. One ED can be enabled by NFV and offer computing services to multiple devices through generating virtual machines, so that various tasks and network functions can be simultaneously performed. On one hand, computation burdens on the SDN controller and transmission backbone can possibly be alleviated by edge computing. On the other hand, the configuration and management costs generated by edge computing can be reduced by NFV and SDN [11]. An SDN-enabled edge computing architecture normally contains three layers: edge layer, SDN infrastructure, and northbound applications. Some beneficial fields in edge computing supported by SDN can be summarized as: high-efficiency resolution and control, flexibility and unrestrained innovation, service-oriented implementation, virtual machine deployment, interoperability, adaptability, economical solution and multiplicity of scope [12]. Furthermore, the integration of SDN and edge computing can be regarded as a distributed SDN, in which the corresponding network requirements of edge computing (such as energy consumption, latency and storage) should also be reached.

The Designed Framework

Deployed EDs are likely to lose efficiency due to inter-channel interference. Also, the process of information flow would cause large delays because of the failure of EDs. With the purpose of guaranteeing network sustainability, we put forward a green and sustainable VNE framework for collaborative edge computing in CoTs. Table 1 summarizes the main notations in this article.

An Overview of the GSVNE Framework

Our two-stage framework is illustrated in Fig. 2. In the first stage, we divide WMNs and IoTs into working and backup sub-parts. For the sake of reducing the consumed resources on the backup sub-graphs, we select $k$ EDs for backup-resource provisioning. After obtaining the value of $k$ by the network reliability function, we derive the number of backup EDs for network resource division. Furthermore, the backup graph is constructed before confirming the geographical locations of the backup EDs. By dividing the resources of WMNs and IoTs, the second stage embeds VNs onto the most appropriate working graph by power control before selecting a suitable backup ED. Next, we conduct the green and sustainable embedding operation for VNs on the most qualified WMN or IoT to guarantee the maximum sharing degree (SD) of backup network resources, representing the maximal amount of virtual nodes embedded onto the same backup ED. Because a constant path exists between the bottleneck and another WMN or IoT with the backup ED, link mapping is not necessary. Therefore, the target of the second phase becomes minimizing the quantity of the consumed backup carriers along fixed paths.

Problem Formulation

The WOBAN in our framework has $n$ WMNs and IoTs. The $i$th ($i \in [1, n]$) WMN or IoT has $M_i$ EDs and $P$ wireless routers, where $M$ and $P$ are the sets of EDs connected to the same ONU and deployed at wireless routers, respectively. For one WMN or IoT, the current VN can be expressed by a four-tuple model $(s, w, b, c)$. The sets of virtual nodes mapped on the working ED connecting to ONU and at wireless routers are illustrated by $s$ and $w$, respectively. The sets of consumed radio bandwidth and computing resources for the working EDs are denoted by $b$ and $c$, respectively. The following steps describe how to embed the virtual nodes onto the working graph.

We first choose a working ONU-level ED to find the mapping solution. After that, we select a collection of working EDs at wireless routers to hold the other virtual nodes, so that all the mapped EDs are with more computing resources than the required resources for VN serving. In addition, the required bandwidth of wireless channels for link mapping should be larger than the consumed bandwidth for mapping by network virtualization. Then, the VN model becomes a two-tuple model denoted by $v(s, w, b, c)$, containing the collection of working EDs mapped by the virtual nodes and the consumed computing resources by EDs, respectively. Once the loca-
tions of \( k \) backup EDs are confirmed, the VN model is embedded onto the backup graph \( \Gamma_b \) by selecting the backup ED with the highest SD among the remaining candidates, whose computing resources are no less than \( c \). After the current network is mapped onto \( \Gamma_w \), the maximal SD equals the number of the virtual nodes mapped onto the ED.

We denote the number of embedded VNs processed by WMNs or IoTs by \( N_t \). The binary variables \( n_{i,b} \) and \( n_{i,w} \) identify whether the \( i \)-th VN is embedded onto the working or backup graph of the qualified WMN or IoT \( \Gamma_w \) or \( \Gamma_b \), respectively. For \( i \in [1, N_t] \) and \( b \in [1, n] \), the total assigned transmission power (TP) to EDs is the summation of \( n_{i,b} \) multiple the required TP on the working graph \( TP_{\Gamma_w} \). Similarly, the maximum SD of the backup resource can be calculated by the sum product of \( n_{i,b} \) and the SD of the backup ED on \( \Gamma_b \).

If there exists no local bottleneck, the objective function of our framework is to minimize the allocated TP to EDs for each backup resource SD. Otherwise, we assign another backup ONU-level ED from WMNs or IoTs, so that \((1 + |\emptyset|)\) virtual nodes are within the VN.

We denote \( N_2 \) by the number of VNs successfully processed by the PON backhaul. The number of consumed optical fiber cables can be obtained by rounding the fraction between \( N_2 \) and the number of VNs served by optical fiber cables. The number of consumed backup carriers is rounding the fraction between the number of VNs served by the optical fiber cables multiple \( wb \) and the carrier’s initial bandwidth provisioning on the optical fiber cable (ba). Therefore, the total number of consumed backup carriers to serve VNs is rounding the fraction between \( N_2 \) multiple \( wb \) and \( ba \). The overall optimization objective is to minimize the integration of two parts, that is, the costs of local WMN or IoT embedding and PON-backhaul embedding.

**HEURISTIC ALGORITHM**

Due to the high computational complexity of the integral linear programming problem, we design a heuristic algorithm to solve the formulated optimization problem iteratively. The process of the GSVNE for collaborative edge computing can be expounded by the following steps.

**Step 1**: Initialize the corresponding information of WOBANs and VNs according to the network status and information.

**Step 2**: Confirm the amount of backup EDs at each WMN or IoT according to the reliability function:

\[
P_r(k) = (1 - P_r)^k \times P_r^{(|M|P) - k}
\]

where \( P_r \) denotes the failure probability of ED, and \((|M|P) - k\) is the residual EDs of available wireless routers for working resource units. Deploying more backup EDs can increase the value of \( k \), resulting in a higher network reliability. If \( P_r < 0.5 \), the value of \( k \) should be increased until the ED failure probability decreases below a threshold.

**Step 3**: For \( \forall u \in M, \forall v \in P \), we deploy the residual \((k - |M|)\) backup EDs; satisfying the minimum hop of \( u \) and \( v \) is no more than the max-hop count between the locations of the ONU-level ED and other wireless routers.

**Step 4**: The resources of the initial and residual backup EDs are divided into two parts according to a scaling factor \( f \). If the ED is not deployed at the wireless router, the initial backup resource provisioning is 0.

**Step 5**: Select a qualified WMN or IoT for VN serving. For each WMN or IoT, the vertexes with lower computing capacity than \( c \) on \( \Gamma_w \) and \( \Gamma_b \) are removed together with the corresponding wireless channels. Then, the weight factor of wireless channel \((u, v)\) is updated by decreasing the required weighting factor among all the outgoing channels of vertex \( u \) iteratively. Afterward, we remove the wireless channels whose weighting factors are 0 to lower the assigned TP further. If the SD of the available backup graph reaches the maximum value, the most qualified network can be decided under the condition that the corresponding path is also available.

**Step 6**: If the most qualified WMN or IoT is available in the current VN, the amount of \( N_t \) is added by 1. Otherwise, we choose the backup ED from different WMNs or IoTs by iteratively decreasing the consumed computing resources of backup EDs, and the quantity of \( N_2 \) is added by 1. The VNE process fails if the backup ED is unavailable; then the quantity of unembedded VNs is increased by 1. The VNE process continues until the last VN can be processed successfully. Otherwise, go back to Step 5.

The flowchart of the designed algorithm is illustrated in Fig. 3.

**Performance Evaluation**

We select the map of the UC Davis campus [13] for smart university community evaluation. In this simulation, each WMN or IoT has the same number of EDs for simplification. As illus-
trated in Fig. 4, local EDs together with three EDs and six wireless routers are deployed in the test WOBAN. For the VNs, we set $|\phi| = 2$, $c = 1$, $wb = 1$, and $ba = 6$.

The consumed TP between GSVNE and the benchmark without graph cutting is demonstrated in Fig. 5a. As the horizontal axis of Fig. 5a shows, different values correspond to various numbers of $N_1$. We can observe that the total TP enhances as the number of embedded VNs increases, and the total consumed power of GSVNE is decreased by around 10 percent. The main reason is that the working graph for each WMN or IoT can be simplified by reserving wireless channels with low weights in GSVNE.

Next, we fix $f$ to 0.5, and evaluate the performance of the total number of embedded VNs. When the circumstance $N > N_1 = 60$ happens, the VNs cannot be embedded by the WMNs or IoTs. As shown in Fig. 5b, it is obvious that the number of the served VNs by the benchmark without PON backhaul is always along the border line caused by a local bottleneck. However, the sustainability of our method can be well guaranteed with more VNs because the backup EDs can still be selected via the PON backhaul. The increase trend remains nonlinear after $N = 110$, because network resources of the ONU-level backup EDs are also constrained.

Figure 5c compares the maximum SD of the backup resources among the theoretical upper bound, GSVNE, and the benchmark without PON backhaul. Since a portion of backup resources offered by EDs is wasted without PON backhaul, the SD of the benchmark is constant and equals 30, well corresponding to the number of served VNs in Fig. 5b. The maximum SD in our method does not rapidly approach the border line at first, since the PON backhaul is not indispensable when the number of VNs is less. The maximal SD of GSVNE remains stable after $N = 80$, well illustrating why the performance of GSVNE remains stable after $N = 110$ in Fig. 5b. Finally, Fig. 5d illustrates the consumed amount of backup carriers between our method ($W_{GSVNE}$) and the optimal bound ($W_{\text{bound}}$). The utilized number of backup carriers by GSVNE increases as the total number of VNs increases. We notice that the gained performances between GSVNE and the optimal bound overlap when $N = [60, 100]$, which demonstrates the effectiveness of our method to guarantee network sustainability. After that, $W_{GSVNE}$ is lower than $W_{\text{bound}}$ since the network resources of PON backhaul are constrained.

**RESEARCH CHALLENGES**

There are still many challenges and open issues to enable collaborative edge computing in CoTs. In this section, we discuss some important aspects of them as follows.

**SECURITY AND PRIVACY**

Although emerging services are enabled by edge computing, its novel features challenge the current security and privacy frameworks. The main reasons are:

- Heterogeneous network structures and deployment of edge computing in WOBANs cause conventional authentication and trust schemes to lose efficiency.
- Although novel communication and software technologies are designed to support edge computing, it is challenging to judge the role of an edge server, that is, whether it is an attacker or an eavesdropper.
- Virtualized infrastructure at the edge is vulnerable to privacy invasion or denial of service attacks caused by virtual machine manipulation.

Furthermore, providing secured connections to the massive devices in CoTs with various processing abilities is also significant. On one hand, the decentralized feature of EDs in WMNs or IoTs is vulnerable; on the other hand, D2D-enabled communication is advocated to reduce latency and balance network loading.
tional resource management is promising to fulfill energy-efficiency edge computing. Different from on-demand elastic cloud resources, many devices in WMNs or IoTs are supplied by battery. The massive interaction in CoTs without network optimization would exhaust the device’s energy quickly. Computational offloading from devices in WMNs or IoTs to edge devices is promising, and D2D-enabled mobility-aware offloading has become a hot research topic. One possible offloading solution is to jointly consider a user’s mobility information, computational ability and channel state. For example, computation-intensive data can be offloaded to the edge servers to reduce latency, and large-size data can be handled by D2D communications to improve energy efficiency [14]. However, D2D communications would generate severe interference, which challenges network scalability. Therefore, interference cancelation and adaptive power control should also be involved in edge computing.

**Data Abstraction**

Although data abstraction has been widely discussed in cloud computing, it is still challenging in edge computing. On one hand, a huge number of things communicates with the EDs at network edge; on the other hand, the deployed EDs report data to the gateway periodically, such as smart water control, light control and door locks in smart home 1 in Fig. 1. Then the processed data can be sent to the upper layer for service providing. However, challenges exist:

- The formats of the delivered data from devices are various.
- A trade-off has to be made between the degree of data abstraction and data storage.
- How to extract useful information reported by the various and unreliable edge devices.
- Different applicable operations coexist on edge devices [10].

**Caching at EDs**

Cache placement and content delivery are two main processes for caching. Although various research works have focused on cache updating during the placement process, online caching is still very challenging for EDs. Recently, cache deployment in mobile core networks has been studied, but how to alleviate the constrained bandwidth of the backhaul links is still challenging.

**Conclusion**

In this article, we first emphasized the significance of sustainable edge computing in CoTs, and discussed research directions for edge computing. Next, a novel GSVNE framework was presented.
The performance analysis showed that our framework can embed more VNs and effectively approach the maximal SD of backup resources compared with the existing schemes. Although our method is an alternative for high energy-efficiency and sustainable edge deployment in CoTs, the corresponding avenues are not flat. Finally, we listed future research challenges and open issues.

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